Practice 2: Digital Systems

Boolean Algebra:

A few important identities:

1.
$$A \cdot \overline{A} = 0$$
 $A \cdot A = A = A + A$

2.
$$A + \overline{A} = 1 = 1 + x$$

3.
$$A + \overline{AB} = A + B$$
 $\overline{A} + AB = \overline{A} + B$

4.
$$\overline{AB} = \overline{A} + \overline{B}$$
 $\overline{A+B} = \overline{A} \cdot \overline{B}$

Example 1: Function Minimization

Minimize the following expression:

$$(\overline{A} + AB\overline{C}) + (A + \overline{A}\overline{B}C)(\overline{A(\overline{A} + \overline{B} + C)})$$

Solution:

$$(\overline{A} + AB\overline{C}) + (A + \overline{A}\overline{B}C)(\overline{A(\overline{A} + \overline{B} + C)}) =$$

$$\overline{A(\overline{A} + \overline{B} + C)} = \overline{A} + \overline{\overline{A} + \overline{B} + C} = \overline{A} + AB\overline{C} = \overline{A} + B\overline{C}$$

$$(\overline{A} + B\overline{C}) + (A + \overline{B}C)(\overline{A} + B\overline{C}) =$$

$$(\overline{A} + B\overline{C})(1 + (A + \overline{B}C)) = \overline{A} + B\overline{C}$$

Example 2: Complementary Functions

Find the complementary expression for:

$$f = A + BC + \overline{A}\overline{B}$$

Solution:

$$\overline{f} = \overline{A + BC + \overline{A}\overline{B}} = (DeMorgan) =$$

$$= \overline{A} \cdot \overline{BC} \cdot \overline{\overline{A}\overline{B}} = \overline{A} \cdot (\overline{B} + \overline{C}) \cdot (A + B) =$$

$$\overline{A} \cdot (A + B) = \overline{A}B$$

$$= \overline{A}B \cdot (\overline{B} + \overline{C}) =$$

$$B \cdot (\overline{B} + \overline{C}) = B\overline{C}$$

$$= \overline{A}B\overline{C}$$

Example 3: Universality

Implement the following function with NAND and NOT gates:

$$f = ABC + \overline{A}\overline{C} + D$$

Solution:

An important step in implementation of digital circuits is finding the *inverting* representation of the function. To do this we will apply a *double inverting* operator:

Example 4: Multiplexers

A Multiplexer is a complex gate that receives 2^N inputs, N selectors and propagates one output. Muxes are one of the most important and useful gates in digital design. In this example, we will implement a standard logic function using the universality of multiplexers:

Implement the following function using 4:1 multiplexers:

$$f(A,B,C,D) = \sum (0,1,5,7,13,14)$$

00 01 11 10 00 1 0 0 0 01 1 1 1 0 11 0 1 0 0 10 0 0 1 0

01

11

10

01 11 10

Solution:

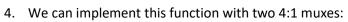
- 1. First, we'll draw the Karnaugh Map of the function:
- A 4:1 Multiplexer receives 2 selectors, so we can make two-variable circles to define the output of each state of the mux. We will choose AB as the selectors for the first mux:
- 3. As we can see, we have arrived at four partial expressions:

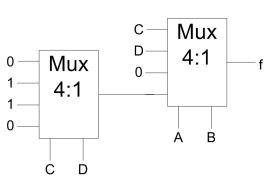
$$f(\overline{A}\overline{B}) = \overline{C}$$

$$f(\overline{A}B) = D$$

$$f(AB) = C \oplus D$$

$$f(A\overline{B}) = 0$$





Bonus Question (if we have time)

- A. Create a leading zeros detector for a four bit vector.
- B. Use the answer to A to create a leading zeros detector for a 16-bit vector.

Bonus 2:

Using 3 Shift Registers and an n-bit adder, design a circuit that receives incrementing integers and outputs their squared value. (1, 4, 9, 16...)